## Santa's Groovy Parallel Helper

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## **Motivation**

 Matt Pedersen's and Jason Hurt's submitted paper to CPA2008



## Critique

- Their JCSP solution did not use two available synchronisation techniques
  - Bucket
    - A component into which one or more processes can fallInto thereby pre-empting themselves, becoming idle, until another process flushes all the processes thereby enabling their re-scheduling
  - Alting Barrier
    - A Barrier enables processes to synchronise such that the set of processes synchronising on the Barrier wait until they have all reached that point in their execution.
    - An Alting Barrier is one that can be used as part of a non-deterministic choice (Alternative)
    - Provides the CSP multi-way synchronisation primitive
- Further simplification by using Groovy Parallel Helper Classes

## **Bucket - methods**

## • fallInto()

- The calling process is pre-empted
- Process becomes idle consuming no processor resource and is associated with the Bucket
- Any number of processes can fall into a bucket
- flush()
  - Must be called by a process that is never pre-empted in a Bucket
  - All the processes associated with the Bucket are rescheduled for execution
    - They may not execute immediately

## **Alting Barrier**

- A possibly dynamic number of processes agree to synchronise on the Alting Barrier
- They do this either
  - absolutely by calling the AltingBarrier's sync() method
  - Process cannot withdraw from the synchronisation
- Or
  - They access the Alting Barrier by means of a guard in an Alternative (ALT)
- Only if the previously agreed number of processes have synchronised or are waiting on an ALT is the Alting Barrier selected as part of a non-deterministic choice

## **Reindeer Synchronisation**

- Alting Barrier comprising
  - Santa Claus
  - Nine Reindeer
- Whenever Santa Claus and the nine reindeer have synchronised on the Alting Barrier
  - Given priority to deliver toys
  - Solely determined when all the reindeer synchronise because Santa checks for this possibility on each iteration
  - Minimal overhead is incurred by Santa
- Implemented as the stable Alting Barrier

## Vestibule

- Contains four groups, each implemented by a Bucket which can each hold up to three elves
- An Elf can tell the Vestibule they need to consult Santa
- The Vestibule then tells the Elf which group (Bucket) to join
- The Elf then fallsInto() the indicated Bucket
- The Elf then waits, idle in the Bucket until it is **flush**ed by the Vestibule.

## **Elf Synchronisation**

- Whenever Santa finishes an elvin consultation he informs the Vestibule
  - The vestibule can then flush() the next group of elves, if any, so they can consult with Santa
- If Santa is idle and a third elf joins a group the Vestibule will flush() the group enabling them to consult with Santa Claus
  - Santa Claus does not have to check to see if there is a waiting group of elves



## Reindeer

def AltingBarrier stable

while (true) {

}

println "Reindeer \${number}: on holiday ... wish you were here, :)" timer.sleep (holidayTime + rng.nextInt(holidayTime)) println "Reindeer \${number}: back from holiday ... ready for work, :(" stable.sync()

harness.write(number)

harnessed.read()

println "Reindeer \${number}: delivering toys . la-di-da-di-da-di-da, :)" returned.read()

println "Reindeer \${number}: all toys delivered ... want a holiday, :(" unharness.read()

## Elf

```
while (true){
        println "Elf ${number}: working, :)"
        timer.sleep ( workingTime + rng.nextInt(workingTime))
        needToConsult.write(1)
        def group = joinGroup.read()
        groups[group].fallInto()
        // idle in Bucket awaiting flush()
        consult.write(number)
        println "Elf ${number}: need to consult Santa, :("
        consulting.read()
        println "Elf ${number}: about these toys ... ???"
        negotiating.write(1)
        consulted.read()
        println "Elf ${number}: OK ... we will build it, bye, :("
```

## **Consult Channel - Elves to Santa**

- Any to One
  - Each elf can write to Santa
- However
  - At any one time only three elves are flushed
  - Hence Santa can expect exactly three communications
  - It does not matter which elf communicates first
  - Provided the other two elf communications are read
- Similarly for the Vestibule channel communications
  - needToConsult (Any2One)
  - joinGroup (One2Any)

## Vestibule – Set Up

```
def flush = new Skip()
def vAlt = new ALT ([needToConsult, consultationOver, flush])
def int index = -1
def int filling = 0
def int removing = 0
def counter = [0, 0, 0, 0]
def NFFD = 0
def OVER = 1
def FLUSH = 2
def preCon = new boolean[3]
preCon[NEED] = true
preCon[OVER] = true
preCon[FLUSH] = false
openForBusiness.read()
```



### Santa – Set Up

def AltingBarrier stable def ChannelInput consult

```
def REINDEER = 0
def ELVES = 1
def rng = new Random()
def timer = new CSTimer()
```

def santaAlt = new ALT([stable, consult])
openForBusiness.write(1)



```
Santa – Elf Choice
case ELVES:
         def id = []
         id[0] = consult.read()
         for (i in 1..2) id[i] = consult.read() // expecting precisely 2 more reads
         println "Santa: ho-ho-ho ... some elves are here!"
         for ( i in 0 .. 2){
                  consulting[id[i]].write(1)
                  println "Santa: hello elf ${id[i]} ..."
         for (i in 0..2) negotiating.read()
         println "Santa: consulting with elves ..."
         timer.sleep ( consultationTime + rng.nextInt(consultationTime))
         println "Santa: OK, all done - thanks!"
         for ( i in 0 .. 2){
                  consulted[id[i]].write(1)
                  println "Santa: goodbye elf ${id[i]} ..."
         consultationOver.write(1)
         break
```

## Result

- Shared Memory (Thread based models)
  - C# 642 lines
  - C 420 lines
  - Java 564 lines
  - Groovy 322 lines
- Distributed Memory
  - MPI 352 lines
- Process Oriented
  - JCSP 315 lines

# Groovy Parallel – 215 lines – 32% reduction over JCSP !!!

