

# Adding Formal Verification to `occam- $\pi$`

Peter H. WELCH<sup>a</sup>, Jan B. PEDERSEN<sup>b</sup>, Fred R.M. BARNES<sup>a</sup>,  
Carl G. RITSON<sup>a</sup> and Neil C.C. BROWN<sup>a</sup>

<sup>a</sup>*School of Computing, University of Kent, UK*

<sup>b</sup>*School of Computer Science, UNLV, USA*

phw@kent.ac.uk, matt@cs.unlv.edu, frmb@kent.ac.uk, cgr@kent.ac.uk, nccb@kent.ac.uk

**Abstract.** This is a proposal for the formal verification of `occam- $\pi$`  programs to be managed entirely within `occam- $\pi$` . The language is extended with qualifiers on types and processes (to indicate relevance for verification and/or execution) and assertions about refinement (including deadlock, livelock and determinism). The compiler abstracts a set of CSP<sub>m</sub> equations and assertions, delegates their analysis to the FDR2 model checker and reports back in terms related to the `occam- $\pi$`  source. The rules for mapping the extended `occam- $\pi$`  to CSP<sub>m</sub> are given. The full range of CSP<sub>m</sub> assertions is accessible, with no knowledge of CSP formalism required by the `occam- $\pi$`  programmer. Programs are proved just by *writing* and *compiling* programs. A case-study analysing a new (and elegant) solution to the *Dining Philosophers* problem is presented. Deadlock-freedom for colleges with *any* number of philosophers is established by verifying an induction argument (the base and induction steps). Finally, following guidelines laid down by Roscoe, the careful use of *model compression* is demonstrated to verify directly the deadlock-freedom of an `occam- $\pi$`  college with 10<sup>2000</sup> philosophers (in around 30 seconds). All we need is a universe large enough to contain the computer on which to run it.

**Keywords.** concurrency, formal verification, model checking, `occam-pi`, CSP, FDR.