Cancellable Servers: a Pattern for Curiosity

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CPA 2012 (University of Abertay, 26th. August, 2012)

Curiosity on Mars

This is a student exercise to design and implement part of the control logic for an autonomous **robot**.control process for a rover vehicle on Mars.

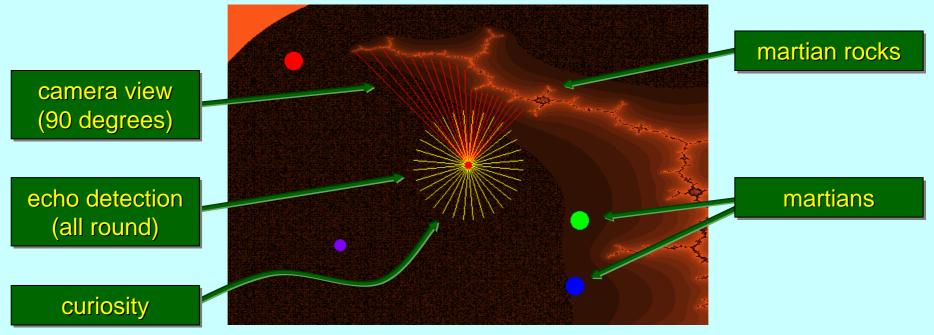






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The controller has to respond to commands from its operator back on Earth, to operate simple actuators (start/stop motors, deploy gadgets) and to monitor and respond appropriately to input from peripherals (motor feedback clicks, raw echo sensor data, processed camera images).



Curiosity on Mars

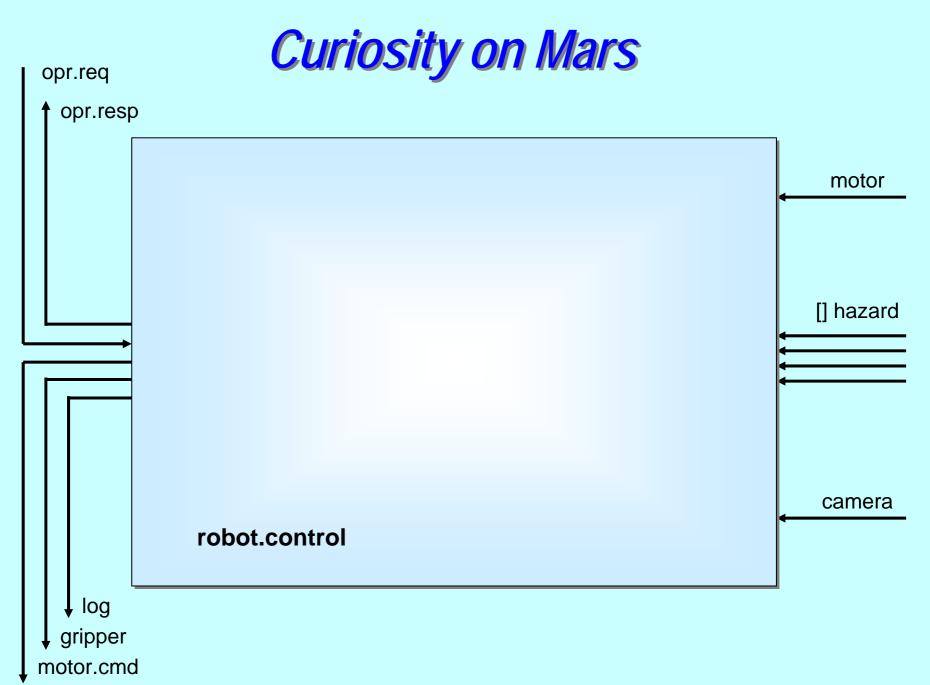
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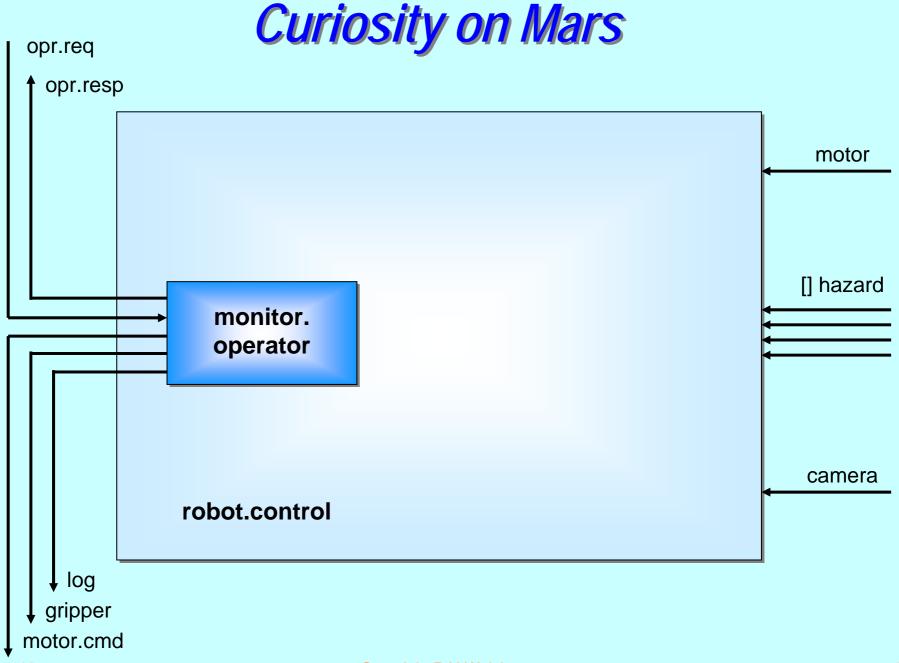
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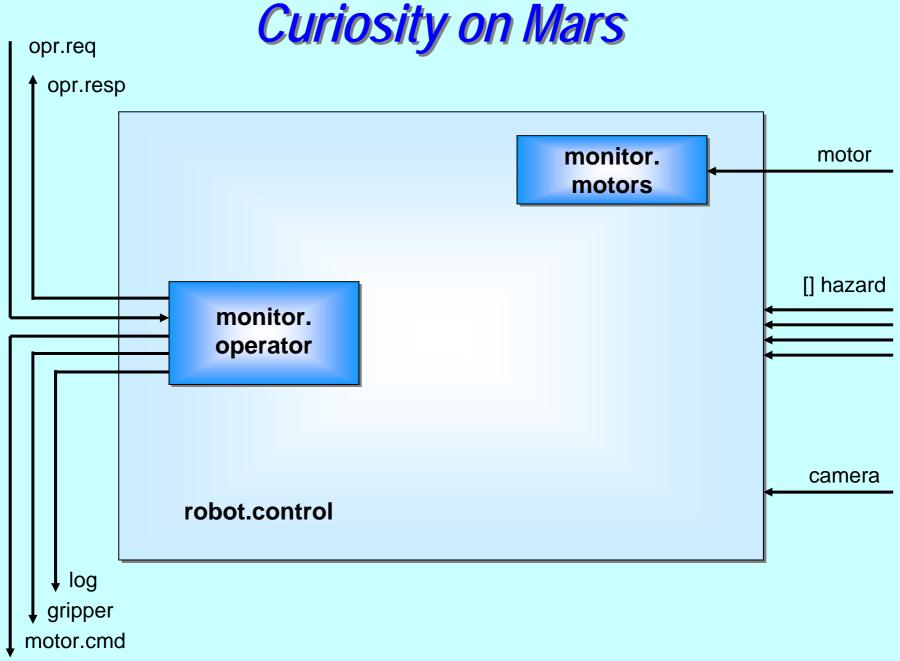
This **could** be implemented by a purely sequential process ... **but that's hard**.

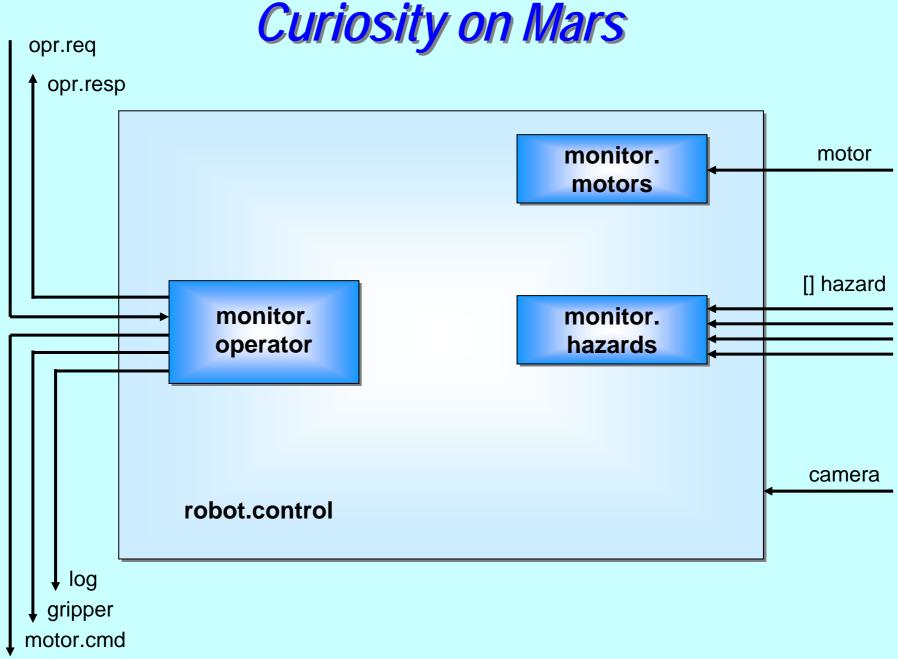
A concurrent implementation is simpler, with a process for each external agent (mission control, motor, echo sounder and camera). These processes are linked and communicate as a *client/server* network.

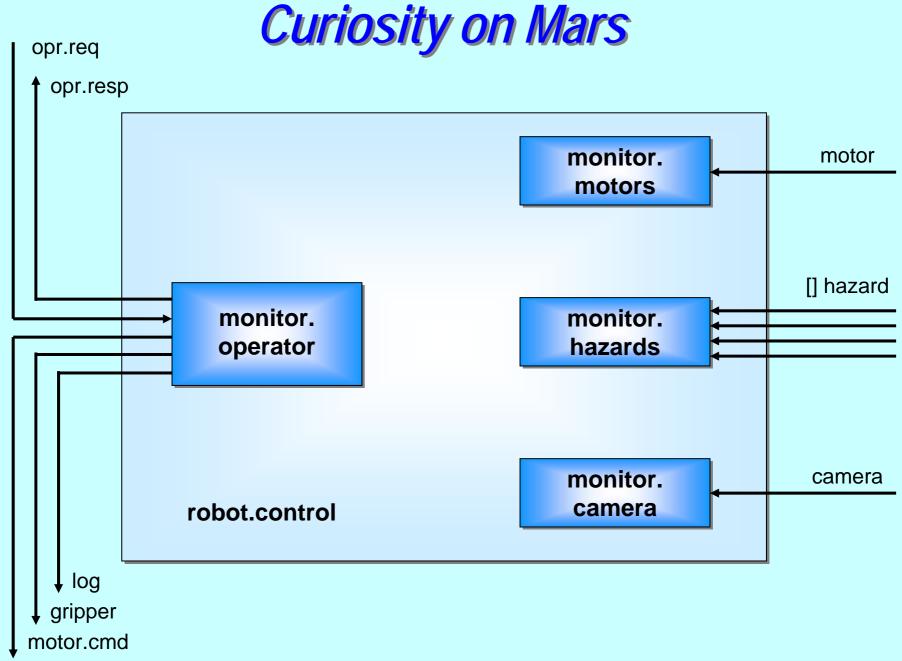
One twist is that three of the servers must be *cancellable* (since two transactions need to run in parallel, with the first to complete causing the cancelation of the other).

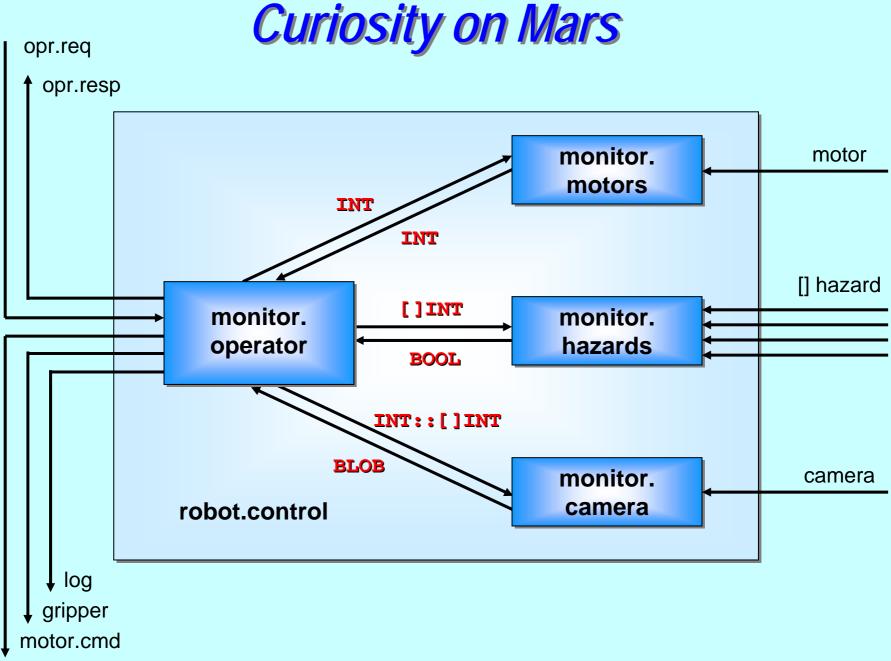


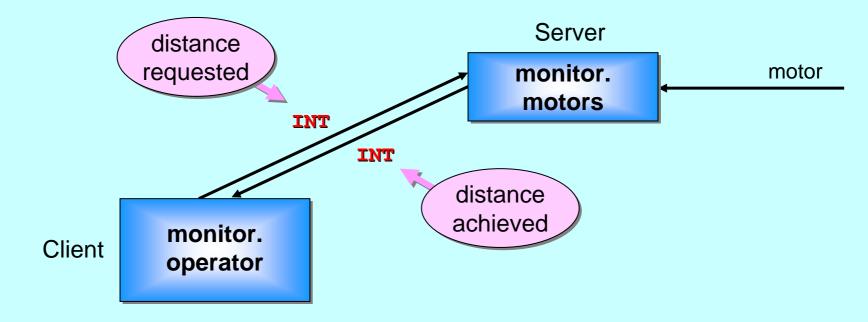


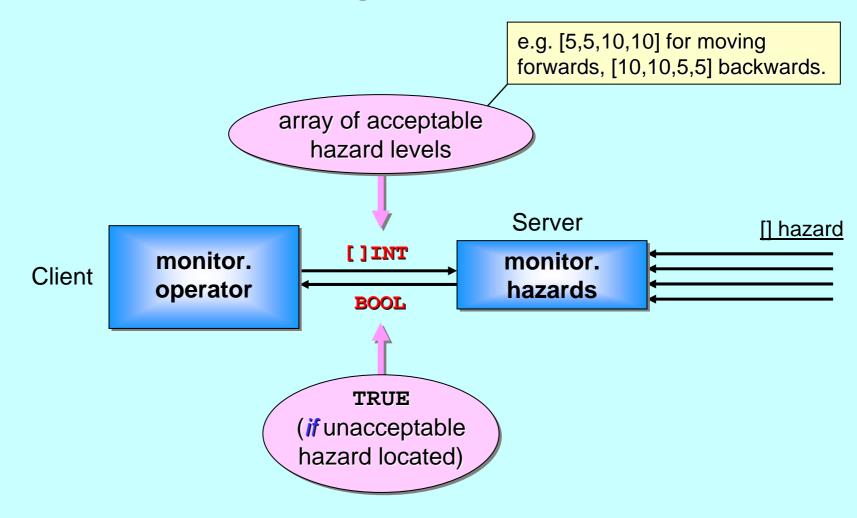


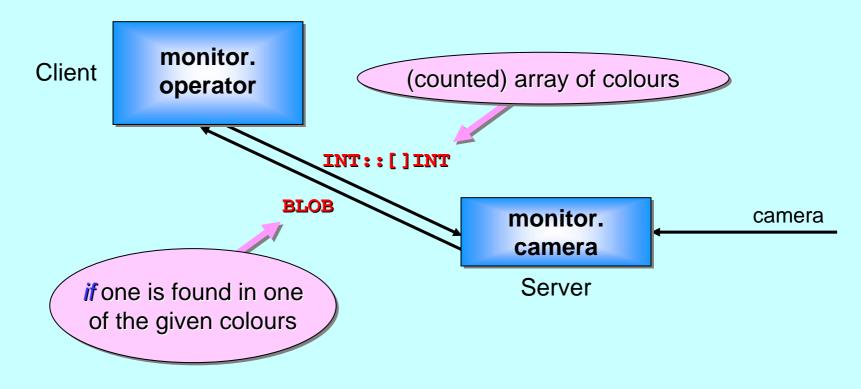


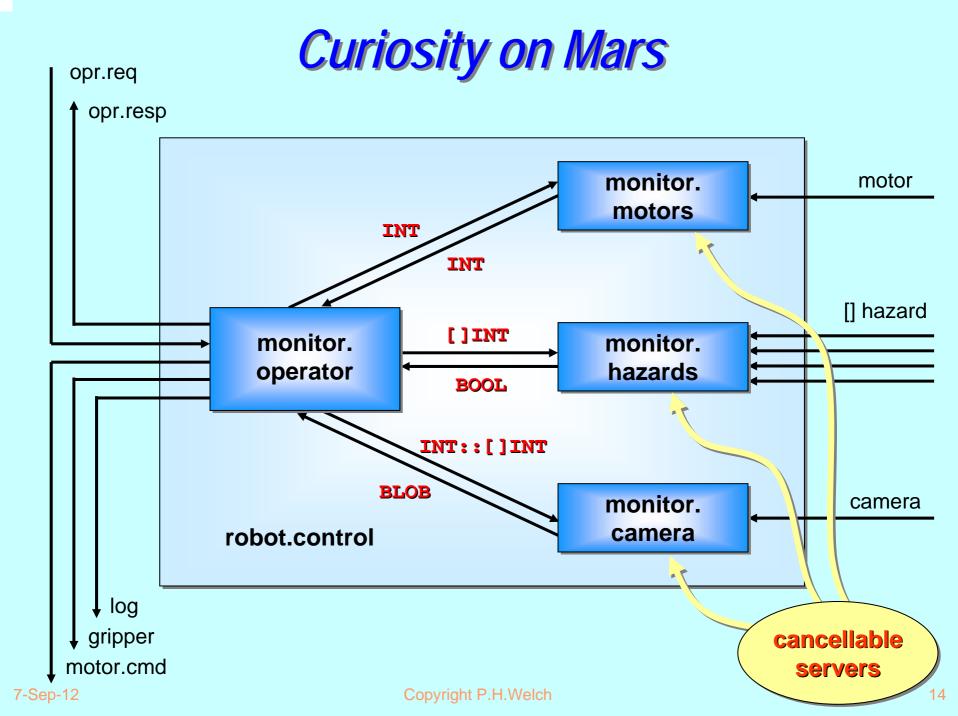


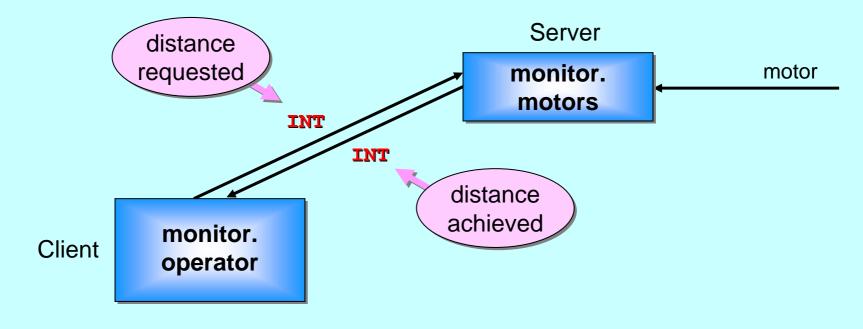






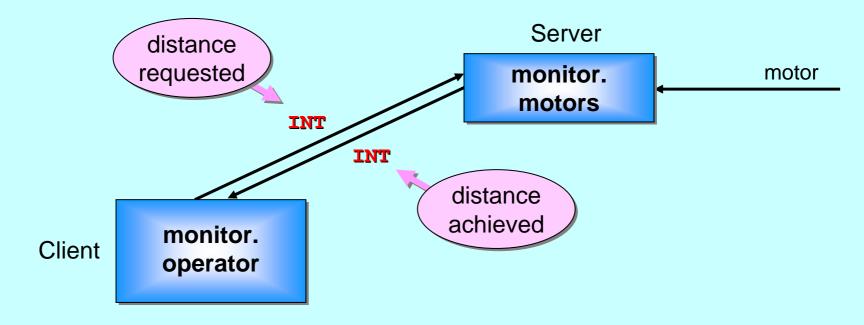






Normal client-server transaction:

- the client sends a request, then waits for an answer.

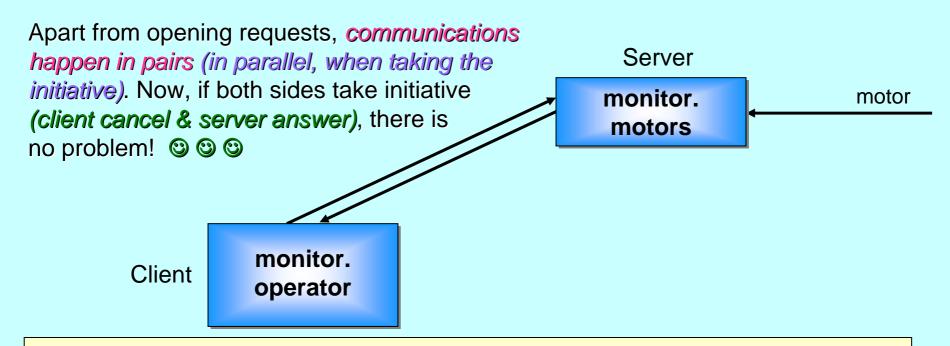


Cancellable client-server transaction:

- the client sends a request, then waits for an answer;
- while waiting for an answer, the client may give up and cancel the request.

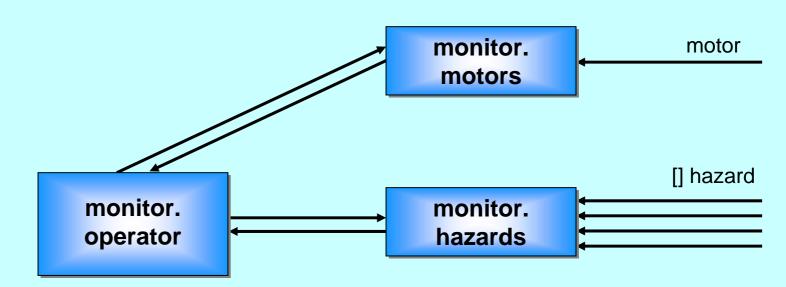
Problem:

- if the client tries to cancel and the server tries to answer, then deadlock!



Solution to cancellable client-server transaction:

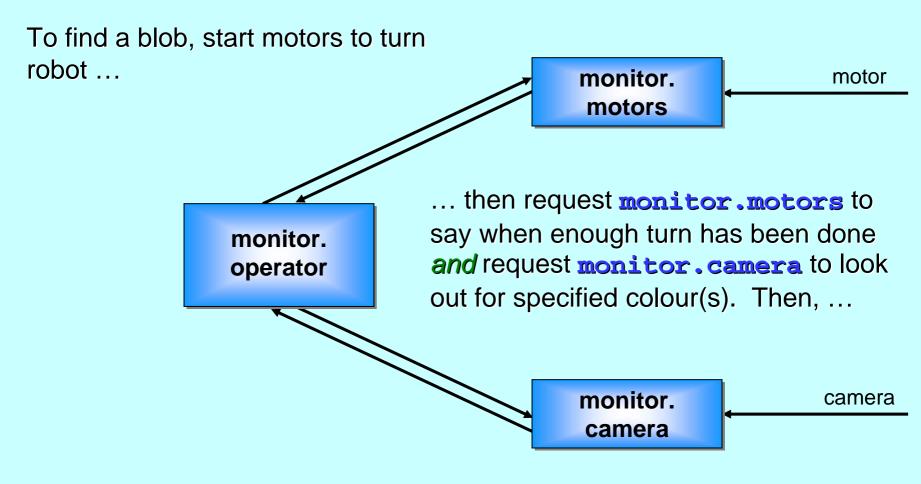
- the client sends a *request*, then waits for an *answer*,
 - If an **answer** is received, client sends an **ack** (confirming receipt);
 - to cancel, the client sends a cancel in parallel with listening for an ack;
 - if an **ack** is received, the request has been cancelled;
 - if an **answer** is received, ignore (server will have seen the **cancel**);
 - to answer, the server sends its answer in parallel with listening for an ack,
 - if an ack is received, server knows client accepted the answer,
 - if a *cancel* is received, server knows the service was cancelled.



To move robot, first check for hazards (normal client-server transaction). If not clear, don't move. Otherwise, start motors to move robot ...

... then request **monitor.motors** to say when enough clicks have been seen **and** request **monitor.hazards** to look out for specified hazards. Listen to both servers for answers. Whoever answers first, cancel the other!

Curiosity on Mars



... listen to both servers for answers. Whoever answers first, cancel the other!

Implement and Verify ...



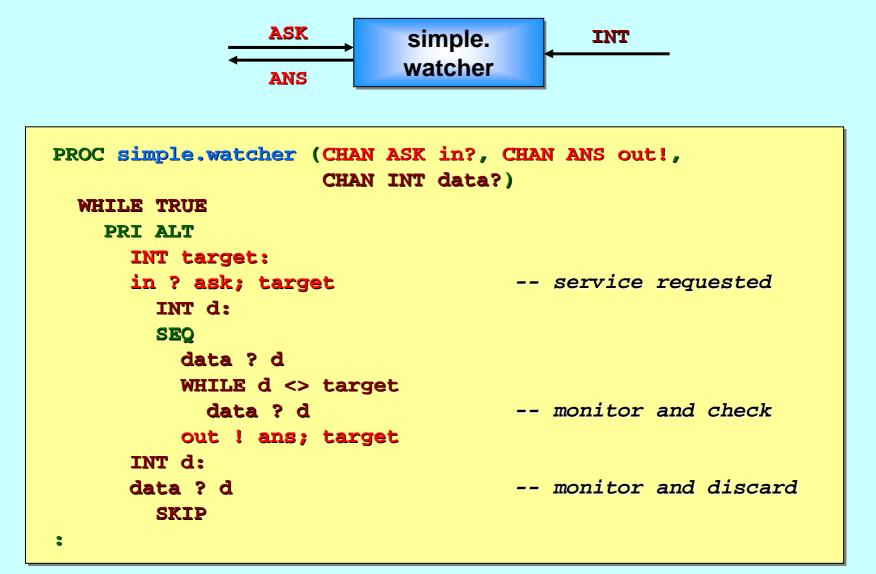
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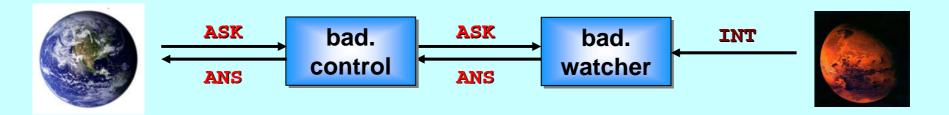


But **simple.watcher** does not deal with a **cancel** request ...

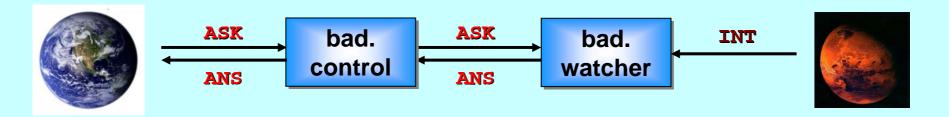
First, let's try it the obvious, but wrong, way ...

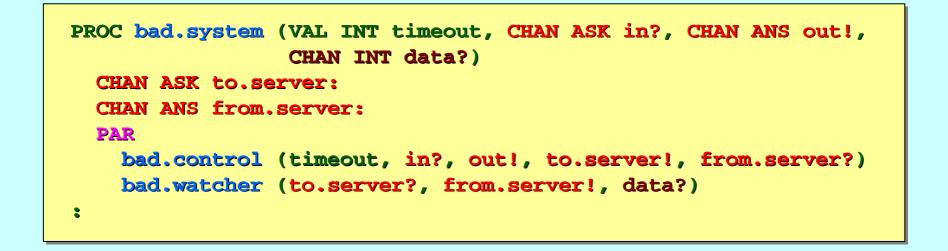


PROC bad.watcher (CHAN ASK in?, CHAN ANS WHILE TRUE	out!, CHAN INT data?)
PRI ALT	
INT target:	
in ? ask; target	service requested
INITIAL BOOL serving IS TRUE:	Service requested
WHILE serving	
PRI ALT	
in ? cancel	service cancelled
serving := FALSE	
INT d:	
data ? d	monitor and check
IF	
d = target	
SEQ	
out ! ans; target	service result
serving := FALSE	
TRUE	
SKIP	
INT d:	
data ? d	monitor and discard
SKIP	
:	

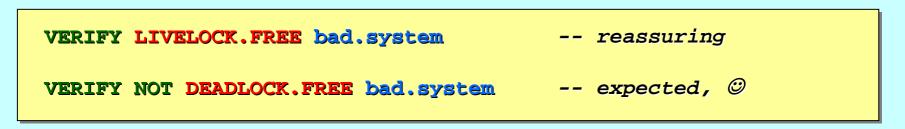


```
PROC bad.control (VAL INT timeout, CHAN ASK in?, CHAN ANS out!,
                 CHAN ASK to.server!, CHAN ANS from.server?)
 WHILE TRUE
   TIMER tim:
    INT t, target:
   SEQ
                                   -- from mission control
     in ? ask; target
     to.server ! ask; target
                                   -- request service
     tim ? t
     ALT
                                    -- service result
       from.server? ans; target
         out ! ans; target
                                  -- to mission control
       tim ? AFTER t PLUS timeout
                                    -- (or PAR)
         SEQ
           to.server ! cancel
                                    -- cancel service
                                    -- to mission control
           out ! ans; -1
```

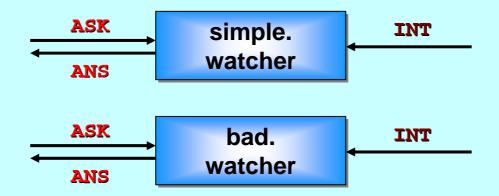




But ...



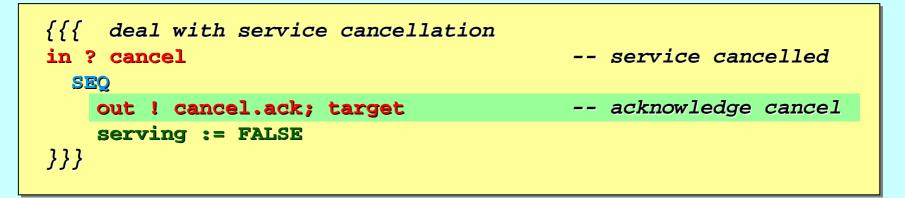
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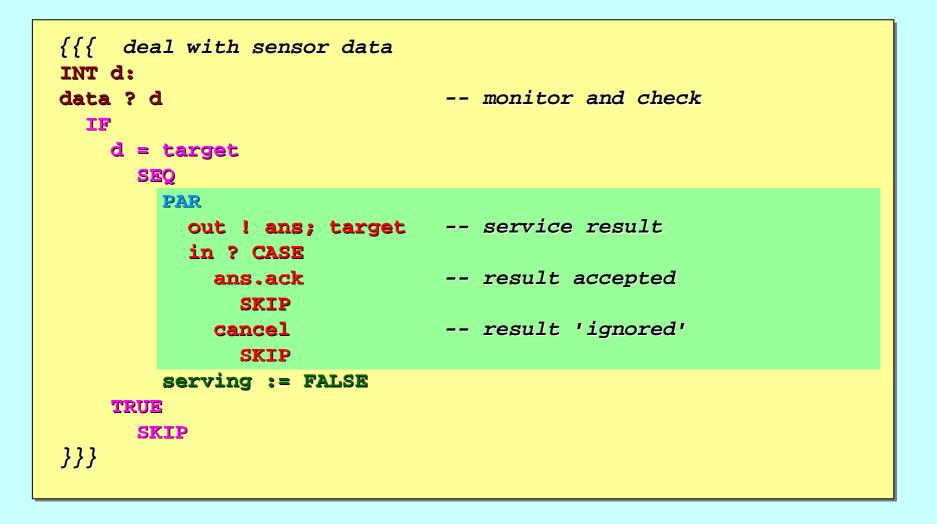


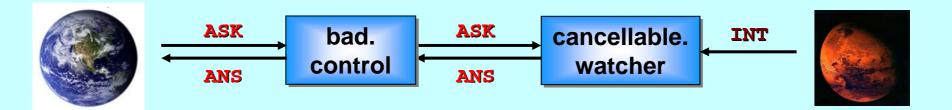
Let's do it right ...

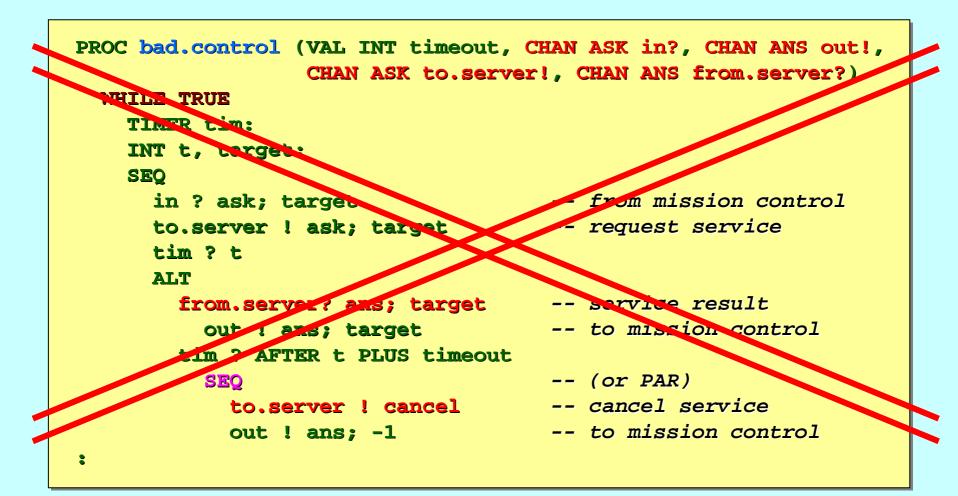


```
PROC cancellable.watcher (CHAN ASK in?, CHAN ANS out!,
                          CHAN INT data?)
  WHILE TRUE
    PRI ALT
      INT target:
      in ? ask; target
                                        -- service requested
        INITIAL BOOL serving IS TRUE:
        WHILE serving
          PRI ALT
            ... deal with service cancellation
            ... deal with sensor data
      INT d:
                                        -- monitor and discard
      data ? d
       SKIP
•
```

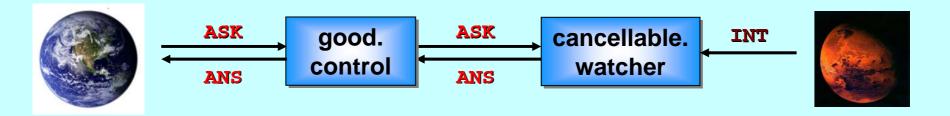


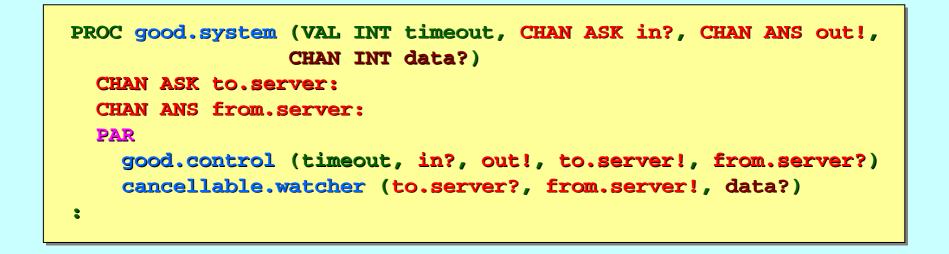






```
PROC good.control (VAL INT timeout, CHAN ASK in?, CHAN ANS out!,
                 CHAN ASK to.server!, CHAN ANS from.server?)
 WHILE TRUE
   TIMER tim:
   INT t, target:
   SEO
                             -- from mission control
     in ? ask; target
     to.server ! ask; target -- forward request
     tim ? t
     ALT
       from.server? ans; target -- service result
                                   -- (or PAR)
         SEO
           to.server ! ans.ack -- acknowledge result
           out ! ans; target -- to mission control
       tim ? AFTER t PLUS timeout
                                   -- (cannot be PAR)
         SEQ
           PAR
             to.server ! cancel
                                   -- cancel service
             from.server ? CASE
                                 -- accept as acknowledge
               ans; target
                 SKTP
               cancel.ack; target -- actual acknowledge
                 SKIP
                                 -- to mission control
           out ! ans; target
2
```





And ...



Curiosity on Mars

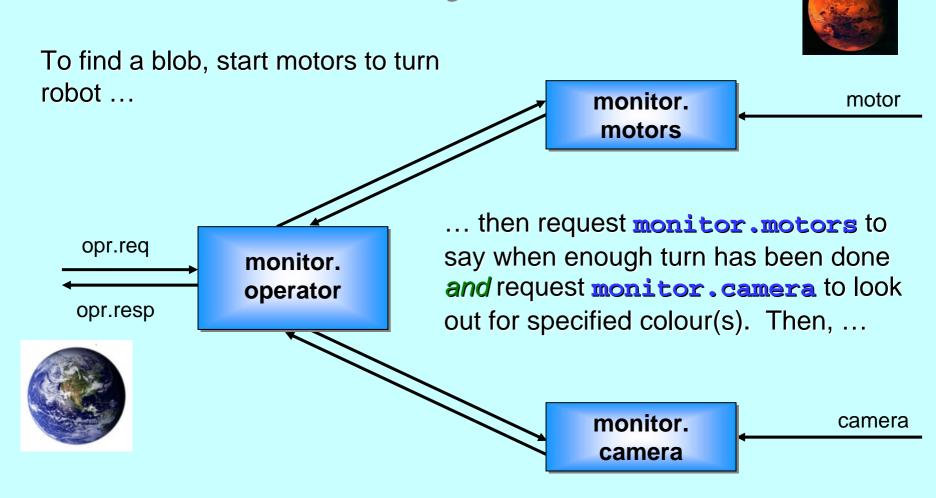
The exercise is to design and implement an autonomous **robot.control** process for a rover vehicle on Mars.

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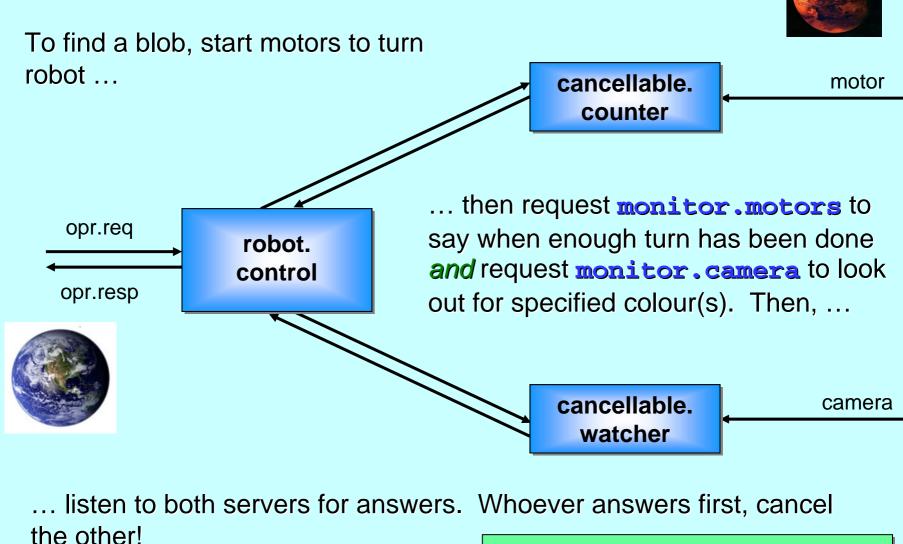
The controller must not deadlock ... or have a sub-system deadlock ...

For Curiosity (or any autonomous vehicle), the verification is not yet sufficient ... we need to verify that *multiple* cancellable servers do not cause problems ...

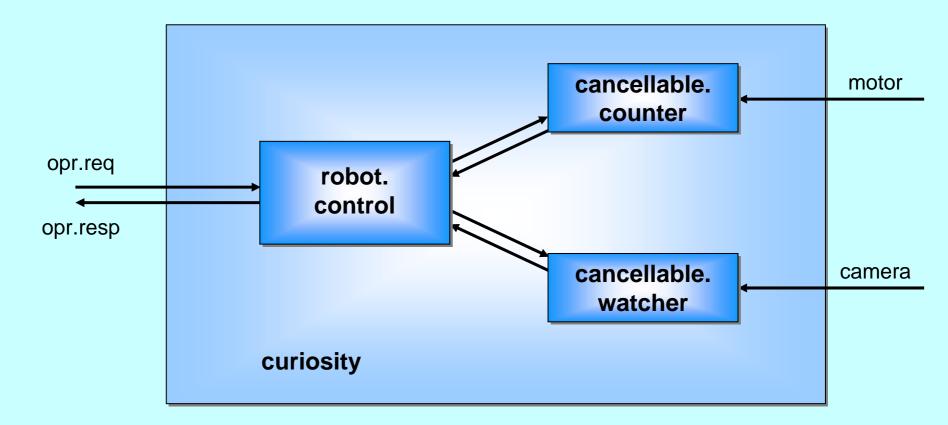
Recall ...



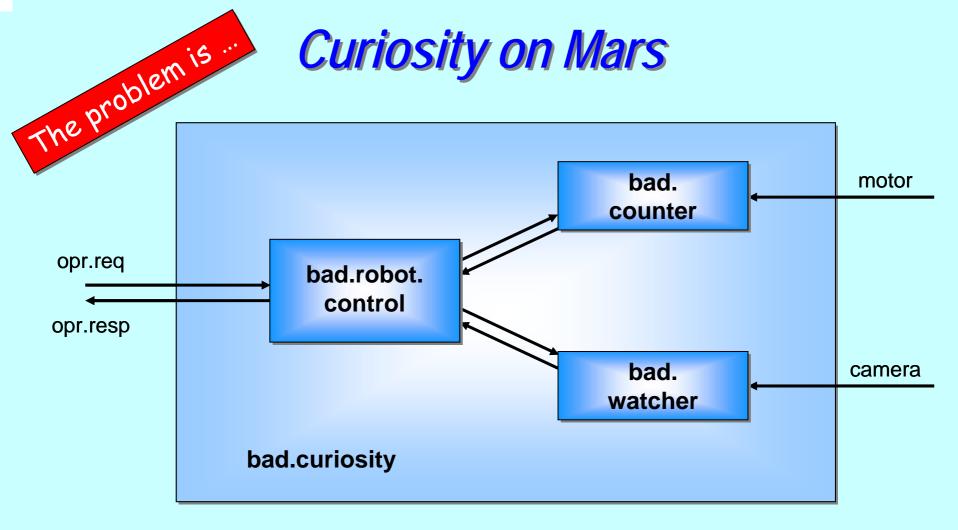
... listen to both servers for answers. Whoever answers first, cancel the other!

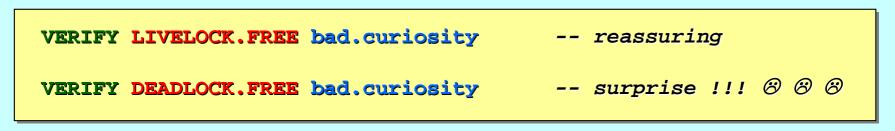


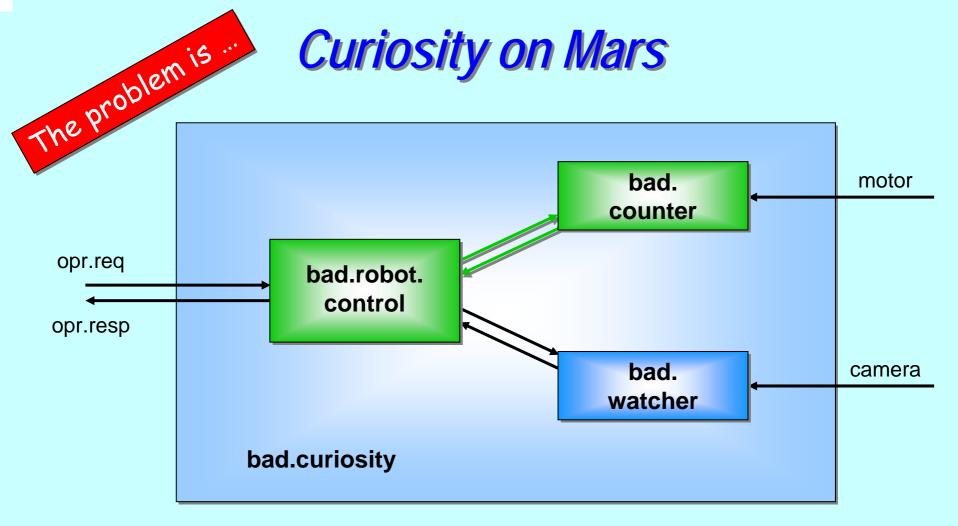
Just changing the names ...



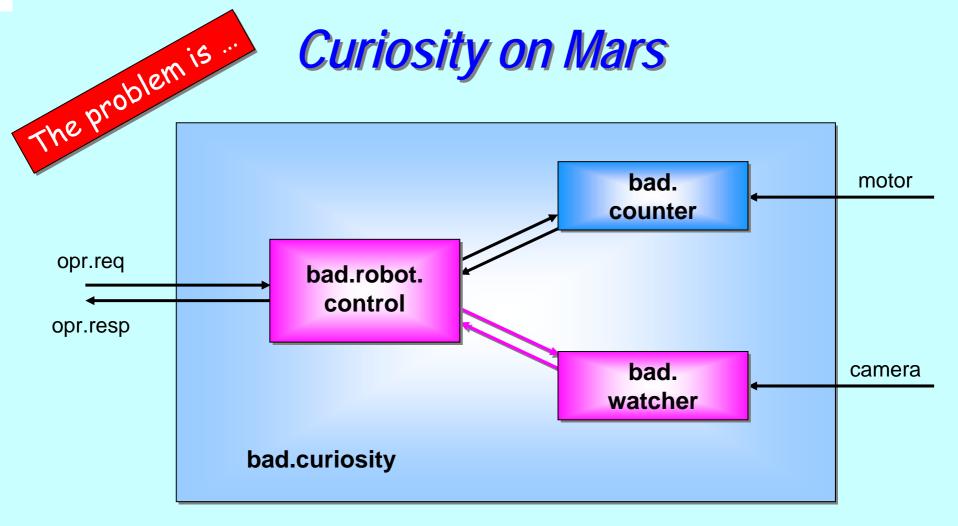




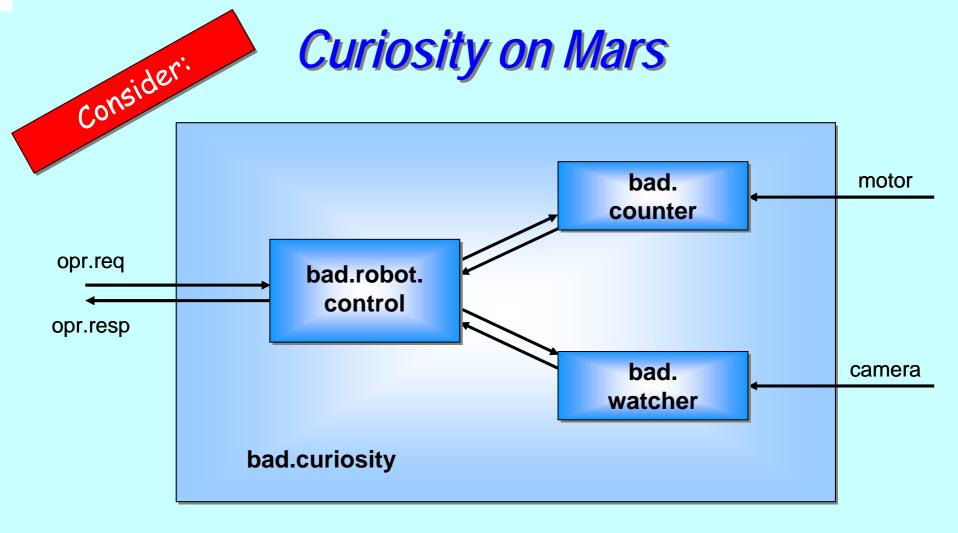


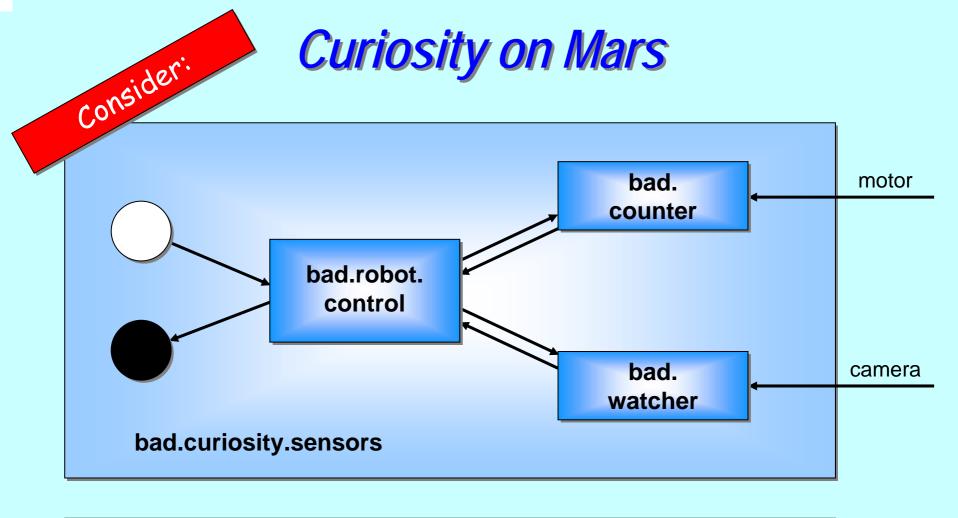


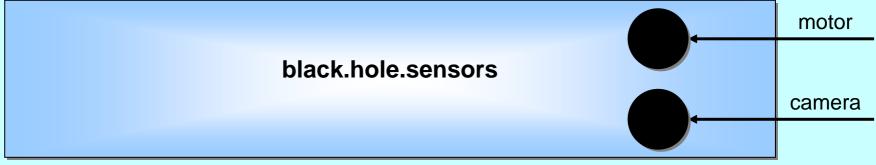
The green sub-system may deadlock, leaving bad.watcher still alive (in its outer loop) accepting and discarding camera data forever. So, the system is not deadlocked!

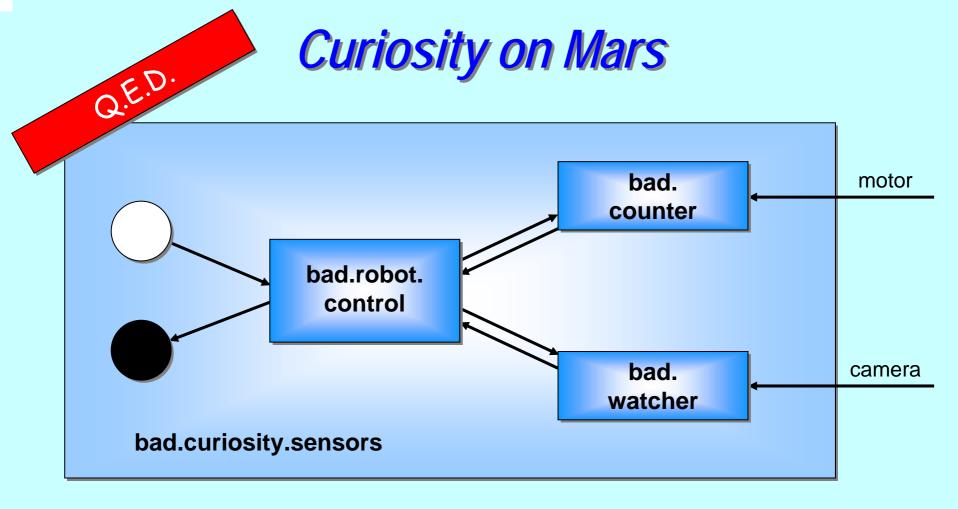


Or the **pink sub-system** may deadlock, leaving **bad.counter** still alive (in its outer loop) accepting and discarding camera data forever. So, the system is not deadlocked!

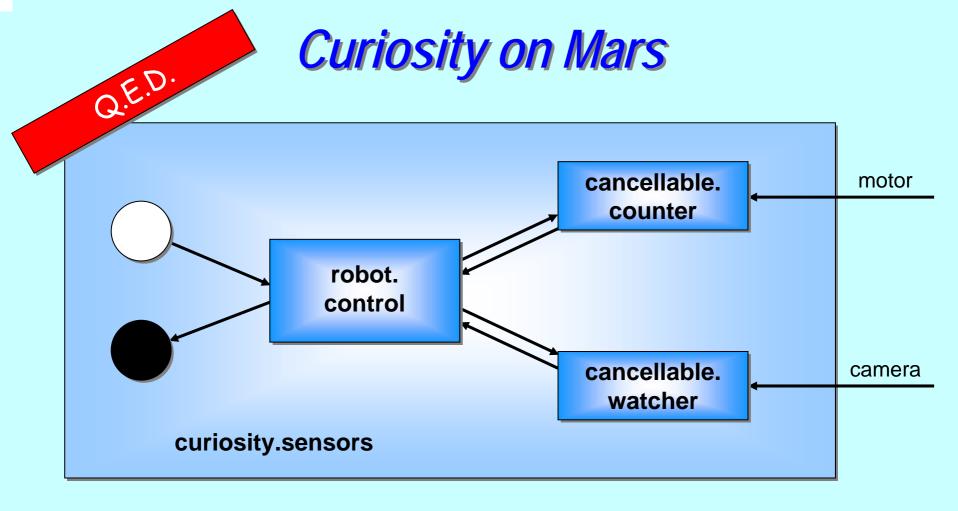




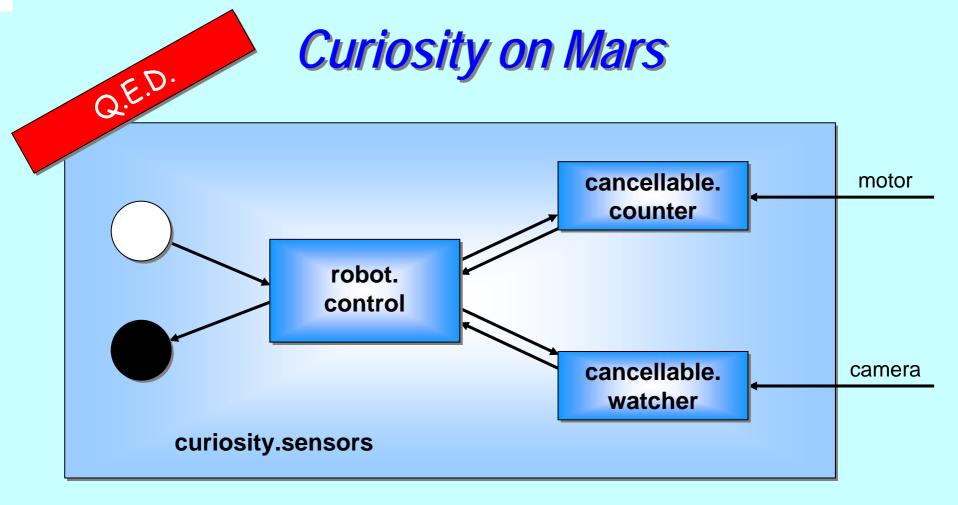








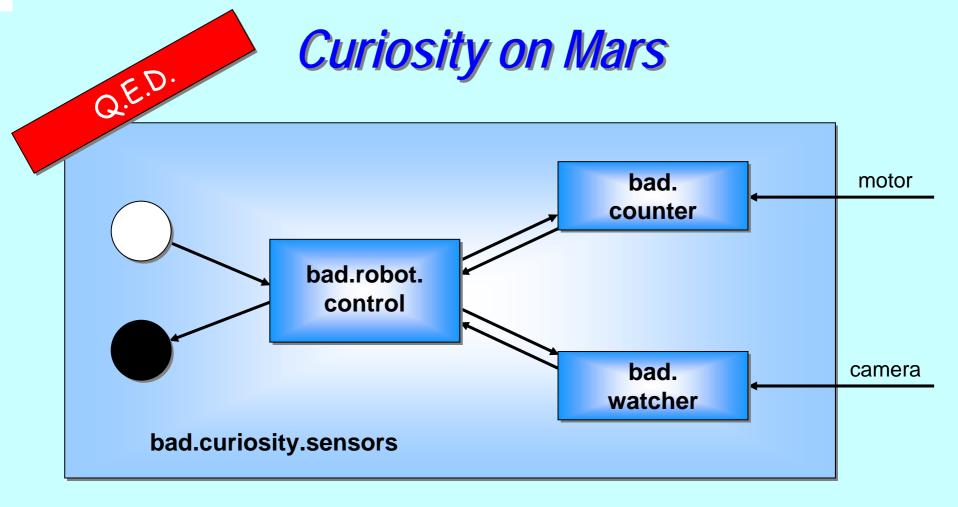






Now, **black.hole.sensors** never refuses motor or camera. Therefore, neither does **curiosity.sensors** (nor **curiosity**).

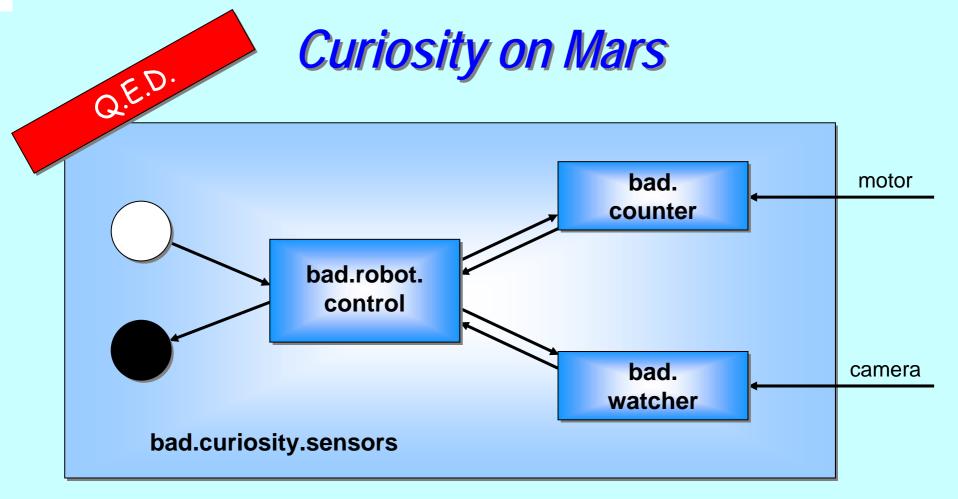
Copyright P.H.Welch





Now, **black.hole.sensors** never refuses motor or camera. Therefore, **bad.curiosity.sensors** (and **bad.curiosity**) does.

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Curiosity on Mars

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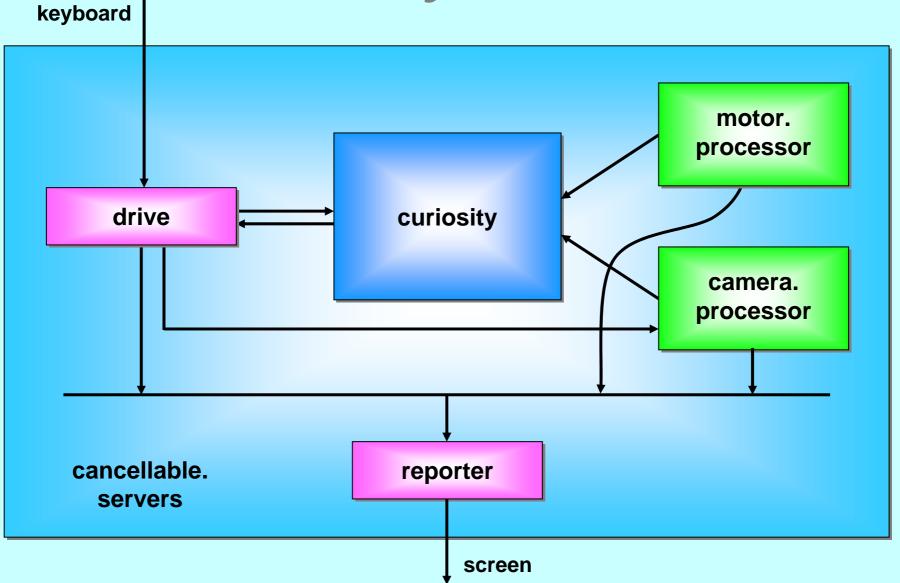
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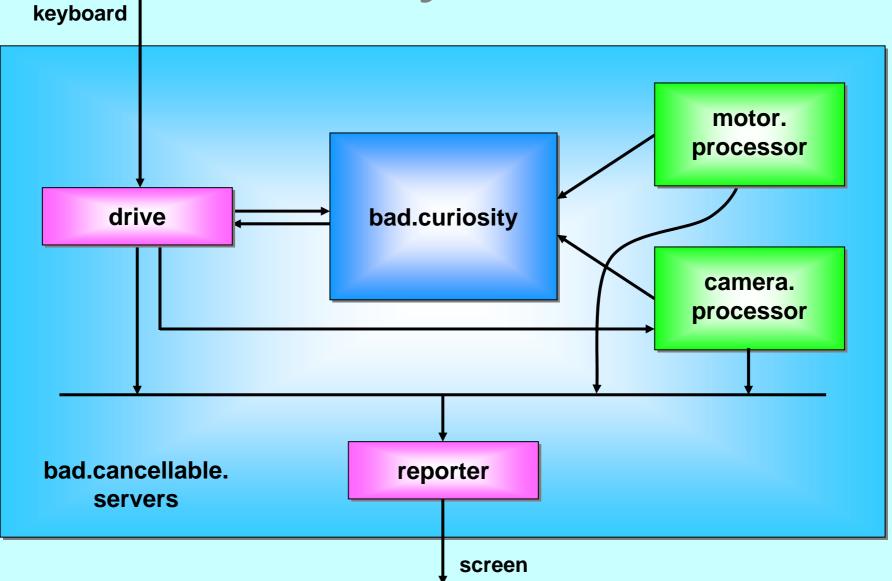


Other Demo ...





Other Demo ...





keyboard motor. processor drive simple.curiosity camera. processor simple. reporter servers screen

Other Demo ...

Curiosity on Mars

Source codes for the system in this presentation is available in 3 forms:

cancellable-servers.op2

occam- π^2 source code (showing generated CSP_m)^{**}

cancellable-servers.csp

CSP_m script (showing **occam**- π^2 source code) – **FDR** ready.

cancellable-servers.occ

Executable **occam**- π source code (with **VERIFY** assertions/**PROC**s commented out) – includes testrig.

* For now, generated by hand ...

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